

Combinatorics Basics

Explained by Examples : Subsets, strings, trees

Master MOSIG : Mathematics for Computer Science

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These notes are only the sketch of the lecture : the aim is to apply the basic counting techniques to the binomial coefficients and establish combinatorial equalities.

References : Concrete Mathematics : A Foundation for Computer Science *Ronald L. Graham, Donald E. Knuth and Oren Patashnik* Addison-Wesley 1989 (chapter 5)



COMPUTER SCIENCE



The four faces of a computer science object

Information representation

Encoding, data, numerical information...

Programming Language

Languages, software engineering...

Algorithmics

Design, proof, complexity,...

Architecture

Processor, networks, operating systems...

Reference : [Les quatre concepts de l'informatique, Didapro, 2011](#) by Gilles Dowek, INRIA/ENS-Saclay

ABSTRACT OBJECTS IN COMPUTER SCIENCE

Symbols, Words, and Texts

- ▶ 011011100101110111...
- ▶ 270c4fe6205c0f43d3163f566534f308
- ▶ grammars, rewriting,...

Sets

- ▶ sets encoding,
- ▶ subsets, partitions

Why counting ?

- ▶ Characteristics of objects
- ▶ Better understanding of the structures
- ▶ Counting = Description Method = Enumeration = Generation

Trees

- ▶ binary trees, binary search trees
- ▶ covering trees
- ▶ tree structures

Ordering

- ▶ permutations
- ▶ sequences
- ▶ partial orders

WELL BALANCED EXPRESSIONS

The problem

$((a + (3 \times (c + 1)) (9 + x) \times ((5 + e) (4 \times 3)))$ is a well-balanced expression ?

and this one

$(5 + a) \times (2 \times 3 \times (5/e)) + (4 - 3 \times (e + 7))$

CHECK_EXPRESSION(S)

Data: S sequence of symbols

Result: True if the expression is well-balanced (else False)

count = 0

for i = 0 to length(S) - 1

 if S[i] == '('

 count = count + 1

 else if S[i] == ')'

 count = count - 1

 if count < 0

 return False

return count == 0

Exercise : Design the algorithm for expressions composed with (), {}, and [], symbols.

ANALYSIS OF THE PROBLEM

Is the algorithm correct ?

- ▶ Formal proof (modify the algorithm to prove it)
- ▶ Check on examples (which ones ?) Is it a proof ?
- ▶ Enumerate all the possible expressions with '(' and ')' and check the correctness. Is it a proof ?
- ▶ Generate a random set arbitrary large of expressions and check. Is it a proof ?

Aim of the activity :

- ▶ Describe the structure, check details, fix notations
- ▶ Explore the small cases exhaustively
- ▶ Establish algebraic structure, make links with other problems

SUBSET ENUMERATION

$\binom{n}{k}$ is the number of ways to choose k elements among n elements



<http://www-history.mcs.st-and.ac.uk/biographies/Pascal.html>

For all integers $0 \leq k \leq n$

$$\binom{n}{k} = \frac{n(n-1)\cdots(n-k+1)}{k!} \quad (1)$$

Prove the equality by a combinatorial argument

Hint : the number of sequences of k different elements among n is $n(n-1)\cdots(n-k+1)$ and the number of orderings of a set of size k is $k!$.

BASIC PROPERTIES

$$\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1} \quad (2)$$

Prove it directly from Equation 1

For all integers $0 \leq k \leq n$

$$\binom{n}{k} = \binom{n}{n-k} \quad (3)$$

Prove it directly from 2

Prove it by a combinatorial argument

Hint : bijection between the set of subsets of size k and ???.

Exercise

Give a combinatorial argument to prove that for all integers $0 \leq k \leq n$:

$$k \binom{n}{k} = n \binom{n-1}{k-1} \quad (4)$$

PASCAL'S TRIANGLE

Recurrence Equation

The binomial coefficients satisfy

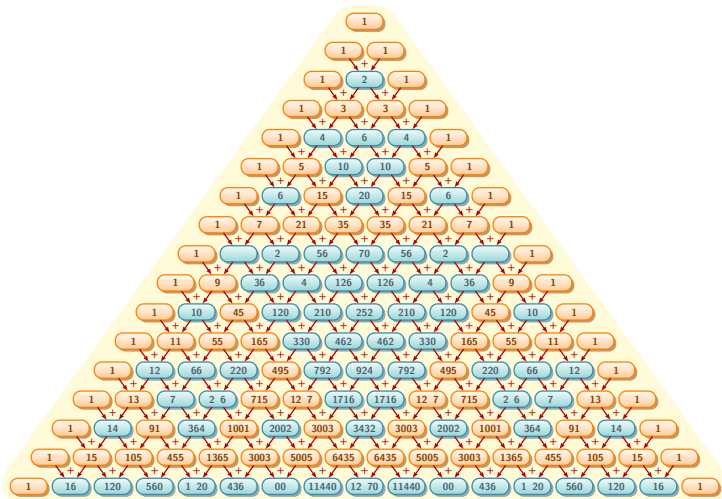
$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} \quad (5)$$

Prove it directly from Equation 1

Prove it by a combinatorial argument

Hint : partition in two parts the set of subsets of size k ; those containing a given element and those not.

PASCAL'S TRIANGLE(2)



Thanks to Tikz/Gaborit

THE BINOMIAL THEOREM

For all integer n and a formal parameter X

$$(1 + X)^n = \sum_{k=0}^n \binom{n}{k} X^k \quad (\text{Newton 1666}) \quad (6)$$

Prove it by a combinatorial argument *Hint : write*

$(1 + X)^n = \underbrace{(1 + X)(1 + X) \cdots (1 + X)}_{n \text{ terms}}$ in each term choose 1 or X , what is the

coefficient of X^k in the result (think vector of n bits).

Exercises

Use a combinatorial argument to prove :

$$\sum_{k=0}^n \binom{n}{k} = 2^n$$

Use the binomial theorem to prove (give also a combinatorial argument)

$$\sum_{k=0}^n \binom{n}{k} = \sum_{k=0}^n \binom{n}{k} = 2^n - 1$$

SUMMATIONS AND DECOMPOSITIONS

The Vandermonde Convolution

For all integers m, n, k

$$\sum_{j=0}^k \binom{m}{j} \binom{n}{k-j} = \binom{m+n}{k} \quad (7)$$

Prove it by a combinatorial argument

Hint : choose k elements in two sets one of size m and the other n .

Exercise

Prove that

$$\sum_{k=0}^n \binom{n}{k}^2 = \binom{2n}{n} \quad (8)$$

Hint : Specify Equation 7

SUMMATIONS AND DECOMPOSITIONS (2)

Upper summation

For all integers $p \leq n$

$$\sum_{k=p}^n \binom{k}{p} = \binom{n+1}{p+1} \quad (9)$$

Exercises

Establish the so classical result

$$\sum_{k=1}^n \binom{k}{1}$$

Compute

$$\sum_{k=2}^n \binom{k}{2} \text{ and deduce the value of } \sum_{k=1}^n k^2$$

THE MAIN RULES IN COMBINATORICS (I)

Bijection Rule

Let A and B be two finite sets if there exists a bijection between A and B then

$$|A| = |B|.$$

Summation Rule

Let A and B be two **disjoint** finite sets then

$$|A \cup B| = |A| + |B|.$$

Moreover if $\{A_1, \dots, A_n\}$ is a partition of A (for all $i \neq j$, $A_i \cap A_j = \emptyset$ and $\bigcup_{i=1}^n A_i = A$)

$$|A| = \sum_{i=1}^n |A_i|.$$

THE MAIN RULES IN COMBINATORICS (II)

Product rule

Let A and B be two finite sets then

$$|A \times B| = |A| \cdot |B|.$$

Inclusion/Exclusion principle

Let A_1, A_2, \dots, A_n be sets

$$|A_1 \cup \dots \cup A_n| = \sum_{k=1}^n (-1)^{k+1} \sum_{S \subset \{1, \dots, n\}, |S|=k} \left| \bigcap_{i \in S} A_i \right|.$$

Exercises

Illustrate these rules by the previous examples, giving the sets on which the rule apply.

THE CENTRAL ROLE OF BIJECTION

Mapping

A **mapping** (function) between X and Y associate to each element x of X a unique element Y

$$\begin{array}{lcl} f : & X & \rightarrow Y \\ & x & \mapsto y \end{array}$$

f is an **injection** iff

$$\forall (x_1, x_2) \in X^2 \quad f(x_1) = f(x_2) \Rightarrow x_1 = x_2$$

f is a **surjection** iff

$$\forall y \in Y \quad \exists x \in X \text{ such that } y = f(x)$$

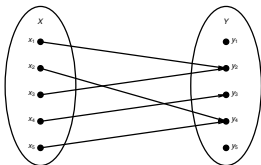
f is a **bijection** iff f is injective **and** surjective

$$\forall y \in Y \quad \exists x \in X \text{ such that } y = f(x) \quad (x \text{ is unique})$$

MAPPINGS AND CARDINALITIES

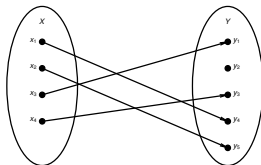
X and Y **FINITE** sets

Typical mapping



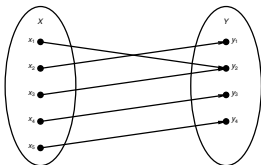
no relation between $|X|$ and $|Y|$

Injective mapping



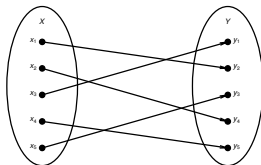
$|X| \leq |Y|$

Surjective mapping



$|X| \geq |Y|$

Bijjective mapping

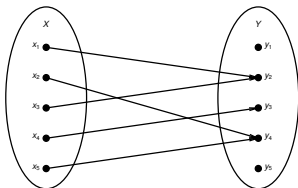


$|X| = |Y|$

What happens when the sets are infinite ?

RECIPROCAL MAPPING

A typical mapping f



$$f^{-1}(y) = \{x \in X, \text{ such that } f(x) = y\}$$

Combinatorial property :

$$\sum_{y \in Y} |f^{-1}(y)| = |X|$$

Exercise :

For all the previous combinatorial proofs construct the corresponding functions.

Inverse Image

subsets of elements of X (equivalence relation on X)

$$f^{-1}(y_1) = \emptyset$$

$$f^{-1}(y_2) = \{x_1, x_3\}$$

$$f^{-1}(y_3) = \{x_4\}$$

$$f^{-1}(y_4) = \{x_2, x_5\}$$

$$f^{-1}(y_5) = \emptyset$$

COUNTING FUNCTIONS (EXERCISES)

Let X and Y finite sets

$$\begin{array}{lcl} f : & X & \rightarrow Y \\ & x & \mapsto y \end{array}$$

- ▶ Compute the total number of such functions f
- ▶ Compute the number of *injective* functions
- ▶ Compute the number of *surjective* functions
- ▶ Compute the number of *bijective* functions

Counting relations

Let X be set, a **relation** \mathcal{R} is a part of $X \times X$.

When X is finite, compute the number of relations on X that are

- ▶ **reflexive** (\mathcal{R} is reflexive iff $\forall x \in X$ we have $x\mathcal{R}x$)
- ▶ **symetric** (\mathcal{R} is symetric iff $\forall(x, y) \in X^2$ we have $x\mathcal{R}y \implies y\mathcal{R}x$)
- ▶ **antisymmetric** (\mathcal{R} is antisymmetric iff $\forall(x, y) \in X^2$ we have $(x\mathcal{R}y \text{ and } y\mathcal{R}x) \implies x = y$)

\mathcal{R} is **transitive** iff $\forall(x, y, z) \in X^3$ we have $(x\mathcal{R}y \text{ and } y\mathcal{R}z) \implies x\mathcal{R}z$

Try to understand why computing the number of transitive relations is hard. [OEIS](#)

DISTRIBUTION PROBLEMS

Context

Place a set of N objects, called *balls*, into a set of M containers, called *urns*.

Basic situations :

- ▶ Labelled balls
- ▶ Labelled urns

More constraints :

- ▶ at least k balls per urn
- ▶ at last k balls per urn
- ▶ number of empty urns
- ▶ ...

EXAMPLE WITH $N = 3$ AND $M = 2$

| | Labelled urns | | Unlabelled urns | |
|------------------|---------------|-------------|-----------------|-------------|
| labelled balls | urn 1 | urn 2 | | |
| | 123 | \emptyset | one urn | the other |
| | 12 | 3 | 123 | \emptyset |
| | 13 | 2 | 12 | 3 |
| | 23 | 1 | 13 | 2 |
| | 1 | 23 | 23 | 1 |
| | 2 | 13 | | |
| | 3 | 12 | | |
| | \emptyset | 123 | | |
| unlabelled balls | urn 1 | urn 2 | one urn | the other |
| | *** | \emptyset | *** | \emptyset |
| | ** | * | ** | * |
| | * | ** | | |
| | \emptyset | *** | | |

Compute the number of configurations in each cell and generalize (if possible).

DERANGEMENT

Definition

A derangement of a set S is a bijection on S without fixed point.
 Number of derangements of n elements d_n (notation $!$).

Inclusion/Exclusion principle

$$\begin{aligned} d_n &= n! - \binom{n}{1}(n-1)! + \binom{n}{2}(n-2)! - \cdots + (-1)^n \binom{n}{n}(n-n)! , \\ &= n! \sum_{i=0}^n \frac{(-1)^i}{i!} \stackrel{n \rightarrow \infty}{\sim} n! \frac{1}{e}. \end{aligned}$$

Recurrence relation

Show by a combinatorial argument that

$$d_n = (n-1)(d_{n-1} + d_{n-2}) = nd_{n-1} + (-1)^n.$$

PROOF OF THE SECOND EQUATION

First we have the first element Thanks [OEIS](#)

| | | | | | | | | | | | | |
|-------|---|---|---|---|---|----|-----|------|-------|--------|----------|-----|
| n | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | ... |
| d_n | 1 | 0 | 1 | 2 | 9 | 44 | 265 | 1854 | 14833 | 133496 | 14684570 | ... |

Suppose that d_n satisfies the recurrence equation $d_n = (n-1)(d_{n-1} + d_{n-2})$ for $n \geq 2$ with $d_0 = 1$ and $d_1 = 0$.

We will prove by recurrence that $d_n = nd_{n-1} + (-1)^n$ with $d_0 = 1$ (E).

1 base case : this is true for $n = 0$ and $n = 1$

2 Suppose that (E) is satisfied for $n-1$

Then $d_{n-1} = (n-1)d_{n-2} + (-1)^{n-1}$, we deduce that $(n-1)d_{n-2} = d_{n-1} - (-1)^{n-1}$.
Injecting that equality in the recurrence equation of d_n

$$\begin{aligned}
 d_n &= (n-1)(d_{n-1} + d_{n-2}) \\
 &= (n-1)d_{n-1} + (n-1)d_{n-2} \\
 &= (n-1)d_{n-1} + d_{n-1} - (-1)^{n-1} \\
 &= nd_{n-1} + (-1)^n
 \end{aligned}$$

3 The base case and the induction is proven, so is the result

FIBONACCI NUMBERS

Recurrence Equation

$$\begin{cases} F_0 = F_1 = 1 \\ F_n = F_{n-1} + F_{n-2} \quad \text{for all } n \geq 2 \end{cases}$$

Interpretation

What kind of situation could be represented by Fibonacci's Numbers ?

Hint : Consider words in $\{0, 1\}^n$

Use a combinatorial argument to prove

$$F_n = F_{n-2} + F_{n-3} + \cdots + F_1 + F_0$$

Hint : Consider the last 1

Imagine other combinatorial equalities

(UNDIRECTED) TREES

A tree $\mathcal{T} = (\mathcal{X}, \mathcal{E})$ is an acyclic connected graph

- ▶ **connected** : for all $x, y \in \mathcal{X}^2$ there is a path from x to y ($x \rightsquigarrow y$)
- ▶ **acyclic** : there are no paths from x to x ($x \rightsquigarrow x$)

Notations

\mathcal{X} set of n nodes

\mathcal{E} set of edges

A **leaf** is a node with exactly one edge and an **internal node** has at least two neighbors.

- ▶ Prove that the maximum number of leaves is $n - 1$ and the minimum 2 (for $n \geq 3$).

An undirected graph \mathcal{T} with n nodes is a tree iff

- 1 \mathcal{T} is acyclic and connected
- 2 \mathcal{T} is acyclic with a maximal number of edges
- 3 \mathcal{T} is connected with a minimal number of edges
- 4 \mathcal{T} is connected with $n - 1$ edges
- 5 \mathcal{T} is acyclic with $n - 1$ edges
- 6 for all couple (x, y) of nodes there is a unique path $x \rightsquigarrow y$ joining the two nodes.

Prove the equivalences (with a minimal number of implications).

CAYLEY'S FORMULA

\mathcal{T}_n the set of all trees with n nodes labelled by the first integers $\{1, 2, \dots, n\}$

T_n the number of such trees.

Phase 1 : small n cases

| n | T_n |
|-----|-------|
| 1 | 1 |
| 2 | 1 |
| 3 | 3 |
| 4 | 16 |
| 5 | 125 |
| ... | ... |

Phase 2 : Intuition of the Formulae

$$T_n = n^{n-2}.$$

Many proofs (see Proofs from the Book).

Approach based on an explicit bijection between the set of trees and the a set of words.

Algorithmic as it associates to each tree a unique word with a coding algorithm.

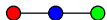
The uniqueness is obtained with a decoding algorithm (H. Prüfer in 1918).



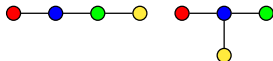
$$= 1$$



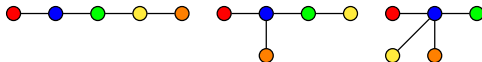
$$2 = 1 = 2^2 \quad 2$$



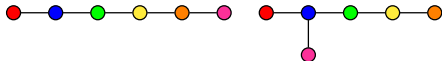
$$3 = \binom{3}{2} = 3 = 3^3 \quad 2$$



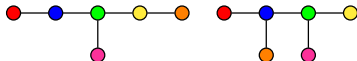
$$4 = \binom{4}{2} \binom{2}{2} \quad 4) = 16 = 4^4 \quad 2$$



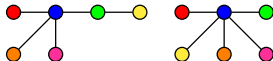
$$5 = \binom{5}{2} \binom{4}{2} \binom{2}{2} \quad 5) \binom{4}{2} \binom{2}{2} \quad 5) \\ = 60 \quad 60 \quad 5 = 125 = 5^5 \quad 2$$



$$6 = \binom{6}{2} \binom{4}{2} \binom{2}{2} \binom{2}{2} \quad 6) \binom{5}{2} \binom{3}{2} \binom{2}{2} \\ \binom{6}{2} \binom{5}{2} \binom{4}{2} \binom{2}{2} \quad 6) \binom{4}{2} \binom{2}{2}$$



$$\binom{6}{2} \binom{5}{2} \binom{2}{2} \quad 6) \\ = 360 \quad 360 \quad 360 \quad 90 \quad 120 \quad 6$$



$$= 1296 = 6^6 \quad 2$$

PRUFER'S CODING ALGORITHM

Phase 3 : double counting

Find a one to one mapping with another set which cardinality is known.

$$\mathcal{T}_n \longleftrightarrow \mathcal{W}_{n-2}$$

\mathcal{W}_{n-2} is the set of words of length $n-2$ over the alphabet $\{1, \dots, n\}$

CODING (T)

Data: A tree T with labelled nodes (all labels are comparable)

Result: A word of $n-2$ labels

$W \leftarrow \{\}$

for $i = 1$ **to** $n-2$ **do**

$x \leftarrow \text{Select_min}(T)$

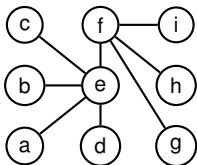
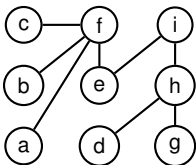
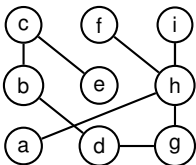
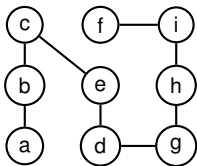
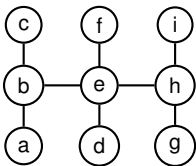
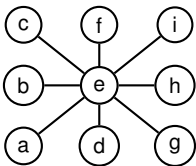
 // x is the leaf with the smallest label

$W \leftarrow W + \text{Father}(x)$

 // **Father** (x) is the unique node connected to the leaf x

$T \leftarrow T \setminus \{x\}$ // remove the leaf x from tree T

LABELLED TREES



PR FER'S DECODING ALGORITHM

DECODING (W)

Data: A word $W = w_1 w_2 \cdots w_{n-2}$ of $n-2$ labels in $\{1, \dots, n\}$

Result: A tree with n nodes labelled from 1 to n

Create n nodes labelled from 1 to n and mark each node by non selected

for $i = 1$ **to** $n-2$ **do**

$x \leftarrow$ **Select_min** (W_i)

 // x is the node with the smallest label not in
 the set $w_1 \cdots w_{i-1}$

 Mark x by selected

 Link x and w_i

Link the last two nodes marked non selected

return T

PR FER'S DECODING ALGORITHM

Examples (10 letters words)

| | | | | | | | | | | |
|----|---|---|---|---|---|---|---|---|---|---|
| 0 | d | i | g | h | a | c | g | c | f | f |
| 1 | e | h | i | e | i | c | a | e | d | d |
| 2 | e | f | g | d | g | g | i | b | c | d |
| 3 | h | h | g | h | c | f | c | c | d | f |
| 4 | i | f | e | c | d | f | a | h | g | f |
| 5 | c | b | e | a | g | i | d | i | a | g |
| 6 | b | g | g | i | b | b | f | i | b | d |
| 7 | e | i | c | c | a | c | f | i | b | d |
| 8 | b | i | d | i | e | e | a | g | d | a |
| 9 | g | c | b | f | c | f | e | f | b | f |
| 10 | b | h | i | a | b | e | b | e | c | h |
| 11 | d | e | h | g | f | f | f | b | e | g |
| 12 | b | h | i | e | a | d | d | g | h | f |
| 13 | g | a | b | h | a | a | g | h | i | i |
| 14 | d | h | d | e | i | i | b | f | b | a |
| 15 | h | e | c | a | b | a | b | c | h | d |
| 16 | i | e | g | i | d | i | e | e | b | g |
| 17 | d | g | i | b | e | h | c | e | i | f |
| 18 | c | h | a | b | e | f | g | b | h | i |
| 19 | h | a | f | b | d | h | c | d | h | g |

Questions

- ▶ Prove the bijection
- ▶ Compute the complexity of coding and decoding
- ▶ What kind of data structure could be useful ?
- ▶ How degrees are expressed in the coding word ?

Extension : is it possible to build a tree from a list of degrees ?

JOYAL'S BIJECTION

What set of objects has cardinality n^n ?

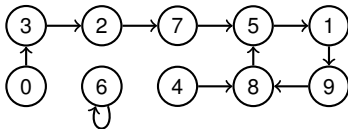
Number of mappings from \mathcal{X} on \mathcal{X} ,

(number of words of size n on an alphabet of size n)

A mapping f

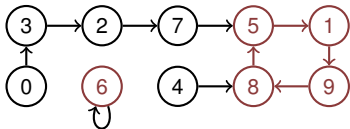
| | | | | | | | | | | |
|--------|---|---|---|---|---|---|---|---|---|---|
| x | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| $f(x)$ | 3 | 9 | 7 | 2 | 8 | 1 | 6 | 5 | 5 | 8 |

Graph associated to mapping f



JOYAL'S BIJECTION

Cycles

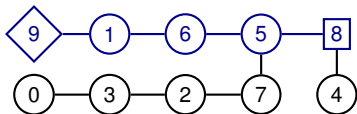


- ▶ Each node has an outdegree = 1
- ▶ Decomposition in **cycles** and **transient** nodes

| | | | | | | | | | | |
|--------|---|---|---|---|---|---|---|---|---|---|
| x | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| $f(x)$ | 3 | 9 | 7 | 2 | 8 | 1 | 6 | 5 | 5 | 8 |

- ▶ Extract the bijective part
- ▶ build a line with the ordered bijective part

Build the tree



- ▶ Fix the line between diamond (image of the smallest) and rectangle (image of the greatest)
- ▶ Connect the transients and remove arrows

Design the reciprocal algorithm

GENERATING FUNCTION

Newton's Binomial Theorem

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} x^k$$

One to one correspondance

$$(1+x)^n \longleftrightarrow \binom{n}{0}, \binom{n}{1}, \dots, \binom{n}{n}$$

Generating Function (Power Series)

Sequence $a = \{a_0, a_1, \dots, a_n, \dots\}$

$$G_a(x) \stackrel{\text{def}}{=} \sum_{n=0}^{+\infty} a_n x^n$$

(formal series, it is not necessary to ensure convergence)

GENERATING FUNCTION (2)

A bijection

Derivation operator

$$G_a(x) = \sum_{n=0}^{+\infty} a_n x^n$$

$$G'_a(x) = \sum_{n=1}^{+\infty} n \cdot a_n x^{n-1}$$

$$G''_a(x) = \sum_{n=2}^{+\infty} n \cdot (n-1) a_n x^{n-2}$$

...

$$G_a^{(k)}(x) = \sum_{n=k}^{+\infty} n \cdot (n-1) \cdots (n-k+1) a_n x^{n-k}$$

...

$$G_a(0) = a_0, \quad \frac{G'_a(0)}{1} = a_1, \quad \frac{G''_a(0)}{2} = a_2, \dots, \quad \frac{G_a^{(k)}(0)}{k} = a_k, \dots$$

BASIC GENERATING FUNCTIONS

| Sequence | \longleftrightarrow Generating function |
|---|---|
| $1, 1, 1, \dots, 1, \dots$ | $\frac{1}{1-x}$ |
| $0, 1, 2, 3, \dots, n, \dots$ | $\frac{x}{(1-x)^2}$ |
| $0, 0, 1, 3, 6, 10, \dots, \binom{n}{2}, \dots$ | $\frac{x^2}{(1-x)^3}$ |
| $1, c, c^2, \dots, c^n, \dots$ | $\frac{1}{1-cx}$ |
| $1, 0, 1, 0, \dots$ | $\frac{1}{1-x^2}$ |
| $\frac{1}{0}, \frac{1}{1}, \frac{1}{2}, \frac{1}{3}, \dots, \frac{1}{n}, \dots$ | e^x |
| $0, \frac{1}{1}, \frac{1}{2}, \frac{1}{3}, \dots, \frac{1}{n}, \dots$ | $\log \frac{1}{1-x}$ |

GENERATING FUNCTIONS : APPLICATIONS

Order one equation

$$a_n = 1 + na_{n-1} \quad n \geq 1$$

$$G_a(x) \quad a_0 = \frac{1}{1-x} \quad 1 + xG'_a(x)$$

Order two equation

$$f_n = f_{n-1} + f_{n-2} \quad n \geq 2$$

Counting objects

Number of ways of choosing a dozen doughnuts when five flavors were available.
 {chocolate, lemon-filled, sugar, glazed, plain}

$$G(x) = \frac{1}{(1-x)^5} = \dots$$

FIBONNACCI'S NUMBERS

Recurrence equation

$f_n = f_{n-1} + f_{n-2}$, for $n \geq 2$, f_0 and f_1 fixed

$G(x)$ generating function of $\{f_n\}$, $G(x) = \sum_n f_n x^n$

$$G(x) - f_0 - f_1 x = x(G(x) - f_0) + x^2 G(x)$$

Decomposition of the generating function

$$G(x) = \frac{f_0 + (f_1 - f_0)x}{1 - x - x^2} = \frac{A}{1 - \varphi x} + \frac{B}{1 - \bar{\varphi} x}$$

with $\varphi = \frac{1+\sqrt{5}}{2}$ and $\bar{\varphi} = \frac{1-\sqrt{5}}{2}$, $1 - x - x^2 = (1 - \varphi x)(1 - \bar{\varphi} x)$

- ▶ Compute A and B and deduce the power expansion of G .
- ▶ Use the power series decomposition $\frac{1}{1-cx} = \sum c^n x^n$.
- ▶ And deduce a closed formula for f_n

GENERATING FUNCTIONS : ALGEBRA

Sequence

 \longleftrightarrow Generating function

$$a_0, a_1, a_2, \dots, a_n, \dots$$

$$G_a(x)$$

$$a_0 + b_0, a_1 + b_1, a_2 + b_2, \dots, a_n + b_n, \dots$$

$$G_a(x) + G_b(x)$$

$$0, a_0, a_1, a_2, \dots, a_{n+1}, \dots$$

$$xG_a(x)$$

$$0 \cdot a_0, 1 \cdot a_1, 2 \cdot a_2, \dots, n a_n, \dots$$

$$xG'_a(x)$$

$$a_0 b_0, a_0 b_1 + a_1 b_0, \dots, a_0 b_n + a_1 b_{n-1} + \dots + a_n b_0, \dots$$

$$G_a(x) \times G_b(x)$$

PIGEONS AND HOLES

Principle

If you have more pigeons than pigeonholes
Then some hole must have at least **two** pigeons

Generalization

If there are n pigeons and t holes, then there will be at least one hole with at least

$$\left\lceil \frac{n}{t} \right\rceil \text{ pigeons}$$

History

Johann Peter Gustav Lejeune Dirichlet (1805-1859)
Principle of socks and drawers



<http://www-history.mcs.st-and.ac.uk/iographies/Dirichlet.html>

SOME EXAMPLES

On integers (from Erdős)

- ▶ Every subset A of $\{1, 2, \dots, 2n\}$ with size $n + 1$ contains at least 2 integers prime together
- ▶ Every subset A of $\{1, 2, \dots, 2n\}$ with size $n + 1$ contains at least 2 integers a and b such that a divide b

On sequences

Consider a sequence of n integers $\{a_1, \dots, a_n\}$.
There is a subsequence $\{a_k, \dots, a_l\}$ such that

$$n \text{ divide } \sum_{i=k}^l a_i$$

IRRATIONAL APPROXIMATION

Friends

Let α be a non-rational number and N a positive integer, then there is a rational $\frac{p}{q}$ satisfying

$$1 \leq q \leq N \text{ and } \left| \frac{p}{q} - \alpha \right| < \frac{1}{qN}$$

Hint : divide $[0, 1[$ in N intervals, and decimal part of $0, \frac{1}{N}, \frac{2}{N}, \dots, \frac{N-1}{N}$

Sums and others

- 1 Choose 10 numbers between 1 and 100 then there exist two disjoint subsets with the same sum.
- 2 For an integer N , there is a multiple of N which is written with only figures 0 and 1

Geometry

- 1 In a convex polyhedra there are two faces with the same number of edges
- 2 Put 5 points inside a equilateral triangle with sides 1. At least two of them are at a distance less than $\frac{1}{2}$
- 3 For 5 point chosen on a square lattice, there are two point such that the middle is also on the lattice

GRAPHS

Friends

Six people

Every two are either friends or strangers

Then there must be a set of 3 mutual friends or 3 mutual strangers

Guess the number

Player 1 : pick a number 1 to 1 Million

Player 2 Can ask Yes/No questions

How many questions do I need to be guaranteed to correctly identify the number ?

Sorting

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