

# Fundamental Computer Science

Denis Trystram  
(inspired by Giorgio Lucarelli)

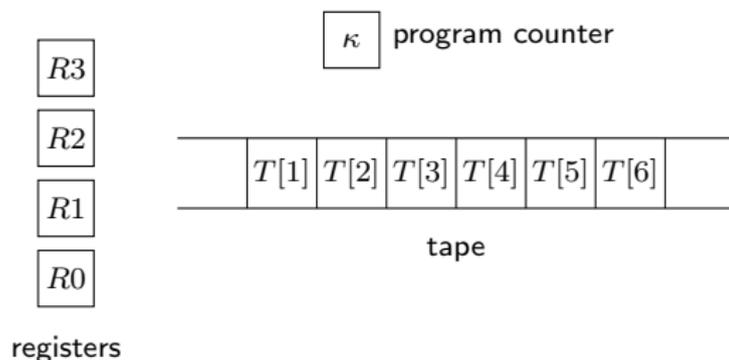
February, 2020

# Random Access Turing Machines

- ▶ Random Access Memory
  - ▶ access any position of the tape in a single step

# Random Access Turing Machines

- ▶ Random Access Memory
  - ▶ access any position of the tape in a single step
- ▶ we also need:
  - ▶ finite number of *registers* → manipulate addresses of the tape
  - ▶ *program counter* → current **instruction** to execute



- ▶ program: a set of instructions

# Random Access Turing Machines: Instructions set

instruction	operand	semantics
read	$j$	$R_0 \leftarrow T[R_j]$
write	$j$	$T[R_j] \leftarrow R_0$
store	$j$	$R_j \leftarrow R_0$
load	$j$	$R_0 \leftarrow R_j$
load	$= c$	$R_0 = c$
add	$j$	$R_0 \leftarrow R_0 + R_j$
add	$= c$	$R_0 \leftarrow R_0 + c$
sub	$j$	$R_0 \leftarrow \max\{R_0 + R_j, 0\}$
sub	$= c$	$R_0 \leftarrow \max\{R_0 + c, 0\}$
half		$R_0 \leftarrow \lfloor \frac{R_0}{2} \rfloor$
jump	$s$	$\kappa \leftarrow s$
jpos	$s$	if $R_0 > 0$ then $\kappa \leftarrow s$
jzero	$s$	if $R_0 = 0$ then $\kappa \leftarrow s$
halt		$\kappa = 0$

► register  $R_0$ : *accumulator*

# Random Access Turing Machines: Formal definition

A Random Access Turing Machine is a pair  $M = (k, \Pi)$ , where

- ▶  $k > 0$  is the finite number of registers, and
- ▶  $\Pi = (\pi_1, \pi_2, \dots, \pi_p)$  is a finite sequence of instructions (program).

# Random Access Turing Machines: Formal definition

A Random Access Turing Machine is a pair  $M = (k, \Pi)$ , where

- ▶  $k > 0$  is the finite number of registers, and
- ▶  $\Pi = (\pi_1, \pi_2, \dots, \pi_p)$  is a finite sequence of instructions (program).

## Notations

- ▶ the last instruction  $\pi_p$  is always a *halt* instruction
- ▶  $(\kappa; R_0, R_1, \dots, R_{k-1}; T)$ : a **configuration**, where
  - ▶  $\kappa$ : program counter
  - ▶  $R_j, 0 \leq j < k$ : the current value of register  $j$
  - ▶  $T$ : the contents of the tape  
(each  $T[j]$  contains a non-negative integer, i.e.  $T[j] \in \mathbb{N}$ )
- ▶ **halted configuration:**  $\kappa = 0$

## Exercise

- ▶ Write a program for a Random Access Turing Machine that multiplies two integers.

Tip: assume that the initial configuration is  $(1; 0, a_1, a_2, 0; \emptyset)$

## Exercise

- ▶ Write a program for a Random Access Turing Machine that multiplies two integers.

Tip: assume that the initial configuration is  $(1; 0, a_1, a_2, 0; \emptyset)$

- 1: **while**  $R_1 > 0$  **do**
- 2:    $R_1 \leftarrow R_1 - 1$
- 3:    $R_3 \leftarrow R_3 + R_2$

# Exercise

- ▶ Write a program for a Random Access Turing Machine that multiplies two integers.

Tip: assume that the initial configuration is  $(1; 0, a_1, a_2, 0; \emptyset)$

- 1: **while**  $R_1 > 0$  **do**
- 2:    $R_1 \leftarrow R_1 - 1$
- 3:    $R_3 \leftarrow R_3 + R_2$

or (all computations should pass through  $R_0$ )

- 1:  $R_0 \leftarrow R_1$
- 2: **while**  $R_0 > 0$  **do**
- 3:    $R_0 \leftarrow R_0 - 1$
- 4:    $R_1 \leftarrow R_0$
- 5:    $R_0 \leftarrow R_3$
- 6:    $R_0 \leftarrow R_0 + R_2$
- 7:    $R_3 \leftarrow R_3$

# Exercise

- ▶ Write a program for a Random Access Turing Machine that multiplies two integers.

Tip: assume that the initial configuration is  $(1; 0, a_1, a_2, 0; \emptyset)$

1: **while**  $R_1 > 0$  **do**

2:    $R_1 \leftarrow R_1 - 1$

3:    $R_3 \leftarrow R_3 + R_2$

1: load 1

2: jzero 9

3: sub =1

4: store 1

5: load 3

6: add 2

7: store 3

8: jump 1

9: halt

or (all computations should pass through  $R_0$ )

1:  $R_0 \leftarrow R_1$

2: **while**  $R_0 > 0$  **do**

3:    $R_0 \leftarrow R_0 - 1$

4:    $R_1 \leftarrow R_0$

5:    $R_0 \leftarrow R_3$

6:    $R_0 \leftarrow R_0 + R_2$

7:    $R_3 \leftarrow R_3$