

# Lecture 3 – Maths for Computer Science

## Solving recurrences and Fibonacci numbers

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Lecture notes MoSIG1

Oct. 2021

## Objective and plan

The purpose of this lecture is to go deeper into recurrence proofs, in particular, bilinear recurrences.

$u_{n+1} = \alpha \cdot u_n + \beta \cdot u_{n-1} + \gamma$  where  $u_0$  and  $u_1$  are given.

## Applications

We already studied such an inductive expression.

### Token Game

$$T(n+1) = T(n) + 2.T(n-1) + 1 \text{ with } T(0) = 1 \text{ and } T(1) = 2$$

## Applications

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### Token Game

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### Fibonacci numbers

The simplest possible bilinear recurrence ( $\alpha = \beta = 1$  and  $\gamma = 0$ ).

$$F(n+1) = F(n) + F(n-1) \text{ with } F(0) = 1 \text{ and } F(1) = 1$$

### Lucas' numbers

Same as Fibonacci with a different seed.

$$L(n+1) = L(n) + L(n-1) \text{ with } L(0) = 1 \text{ and } L(1) = 3$$

### Derangements

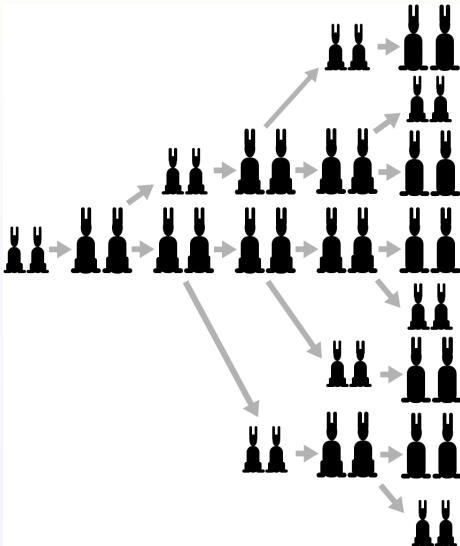
$$d(n+1) = n(d(n-1) + d(n-2)) \text{ with } d(0) = 1 \text{ and } d(1) = 2$$

## Definition of Fibonacci numbers

The original problem has been introduced by Leonardo of Pisa (Fibonacci) in the middle age.

- Fibonacci numbers are the number of pairs of rabbits that can be produced at the successive generations.
- Starting by a single pair of rabbits and assuming that each pair produces a new pair of rabbits at each generation during only two generations.

## Definition (pictorially)



## Definition (more formally)

### Definition:

Given the two numbers  $F(0) = 1$  and  $F(1) = 1$

the Fibonacci numbers are obtained by the following expression:

$$F(n + 1) = F(n) + F(n - 1)$$

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$$F(n + 1) = F(n) + F(n - 1)$$

The first ranks:

n	0	1	2	3	4	5	6	7	8	9	10	...
F(n)	1	1	2	3	5	8	13	21	34	55	89	...



## Combinatorial interpretation

### Proposition

The Fibonacci number  $F(n)$  can be interpreted as the number of length- $n$  binary strings in which each occurrence of a 1 is directly preceded by a 0.

Let  $S_n$  be the set of such strings of length  $n$ .

## Proof

By the previous definition, every binary string  $\omega_n$  ends either with 0 or with 01.

- If  $\omega_n$  ends with 0, then, it has the form  $x0$  where the prefix  $x$  is a binary string of length  $n - 1$ .

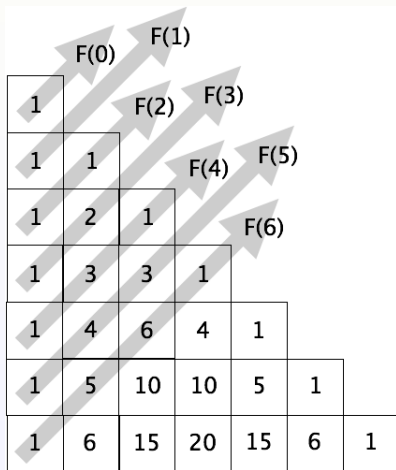
Moreover,  $x$  must belong to  $S_{n-1}$  in order  $\omega_n$  belongs to  $S_n$ . Therefore  $S_n$  contains  $|S_{n-1}|$  strings of this form.

- If  $\omega_n$  ends with 01, then it has the form  $\omega_n = y01$ , where the prefix  $y$  is a binary string of length  $n - 2$ .

Moreover,  $y$  must belong to  $S_{n-2}$  in order for  $\omega_n$  to belong to  $S_n$ , that contains  $|S_{n-2}|$  strings of this form.

$$F(n) = |S_n| = F(n - 1) + F(n - 2)$$

## Link with the Pascal's triangle



## Studying a first property

Proposition:

$$F(n+2) = 1 + \sum_{k=0}^n F(k)$$

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Let check the expression on the first ranks:

$$n = 1,$$

$$F(3) = 1 + F(1) + F(0) = 1 + 1 + 1 = 3$$

$$n = 2,$$

$$F(4) = 1 + F(2) + F(1) + F(0) = 1 + 2 + 1 + 1 = 5$$

$$n = 3,$$

$$F(5) = 1 + F(3) + F(2) + F(1) + F(0) = 1 + 3 + 2 + 1 + 1 = 8$$

...

# Proof

By induction

- The **basis case** (for  $n = 0$ ) is true since  $F(2) = 1 + F(0)$ .
- **Induction step:** Let assume the property holds at rank  $n$  for  $F(n + 2)$  and compute  $F(n + 3)$ :

Apply the definition of Fibonacci numbers:

$$F(n + 3) = F(n + 1) + F(n + 2)$$

Replace the last term by the recurrence hypothesis:

$$F(n + 2) = 1 + \sum_{k=0}^n F(k)$$

Thus,

$$F(n + 3) = F(n + 1) + 1 + \sum_{k=0}^n F(k) = 1 + \sum_{k=0}^{n+1} F(k)$$

## Product of two consecutive Fibonacci numbers

Proposition:

$$F(n).F(n-1) = \sum_{k=0}^{n-1} F(k)^2 \text{ (for } n \geq 1)$$

Let check the expression on the first ranks:

$$n = 2, F(2).F(1) = F(1)^2 + F(0)^2 = 1 + 1 = 2$$

$$n = 3, F(3).F(2) = F(2)^2 + F(1)^2 + F(0)^2 = 4 + 1 + 1 = 6$$

$$n = 4, F(4).F(3) = F(3)^2 + F(2)^2 + F(1)^2 + F(0)^2 = 15$$

$$n = 5, F(5).F(4) = F(4)^2 + F(3)^2 + F(2)^2 + F(1)^2 + F(0)^2 = 40$$

...

## Proof by induction

- The **basis case** (for  $n = 1$ ) is true since  $F(1).F(0) = F(0)^2 = 1$ .
- **Induction step**<sup>1</sup>: Let assume the property holds at rank  $n$  and compute  $F(n + 1).F(n)$ :  
 Apply the definition of  $F(n + 1)$ :  

$$F(n + 1).F(n) = (F(n) + F(n - 1)).F(n)$$

$$= F(n)^2 + F(n).F(n - 1)$$
 Apply now the induction hypothesis to this last term:  

$$F(n + 1).F(n) = F(n)^2 + \sum_{k=0}^{n-1} F(k)^2 = \sum_{k=0}^n F(k)^2$$

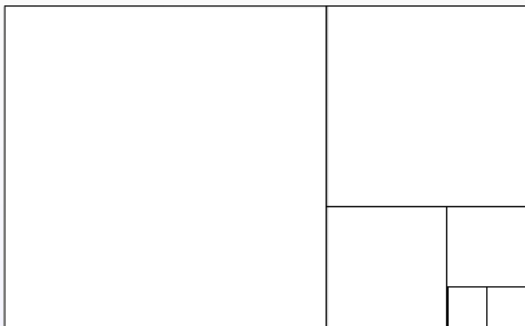
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<sup>1</sup>exactly the same scheme as before!



## An alternative proof by recurrence

The relation can be proved very easily by the geometric argument shown below



## Another property dealing with squares

Proposition:

$$F(n+2)^2 = 4.F(n).F(n+1) + F(n-1)^2 \text{ for } n \geq 2.$$

Let check the expression on the first ranks:

$$n = 1, F(3)^2 = 3^2 = 4.F(1).F(2) + F(0)^2 = 8 + 1 = 9$$

$$n = 2, F(4)^2 = 5^2 = 4.F(2).F(3) + F(1)^2 = 24 + 1 = 25$$

$$n = 3, F(5)^2 = 8^2 = 4.F(3).F(4) + F(2)^2 = 60 + 4 = 64$$

$$n = 4, F(6)^2 = 13^2 = 4.F(4).F(5) + F(3)^2 = 160 + 9 = 169$$

...

## Analytic proof

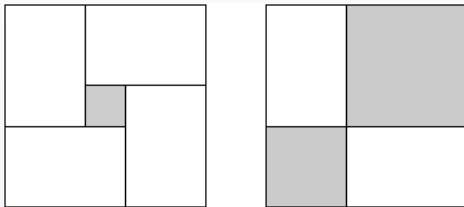
Use the definition of the Fibonacci numbers and expand:

$$\begin{aligned}F(n+2)^2 &= (F(n+1) + F(n))^2 \\&= F(n+1)^2 + 2.F(n+1).F(n) + F_n^2 \\&= 4.F(n+1).F(n) - 2.F(n+1).F(n) + F(n+1)^2 + F(n)^2 \\&= 4.F(n+1).F(n) + (F(n+1) - F(n))^2\end{aligned}$$

Again, using the definition of  $F(n+1)$  into the square, we get the expected result:

$$F(n+2)^2 = 4.F(n+1).F(n) + F(n-1)^2$$

## Graphical proof



## Cassini's identity

Proposition:

$$F(n-1).F(n+1) = F(n)^2 + (-1)^{n+1} \text{ for } n \geq 1$$

Let check the expression on the first ranks:

$$n = 1, F(0).F(2) = F(1)^2 + 1 = 2$$

$$n = 2, F(1).F(3) = F(2)^2 - 1 = 3$$

$$n = 3, F(2).F(4) = F(3)^2 + 1 = 10$$

$$n = 4, F(3).F(5) = F(4)^2 - 1 = 24$$

...

## Proof (by induction)

- The **basis case**  $n = 1$  holds since  $F(0).F(2) = F(1)^2 + 1 = 2$ .
- The **induction step** is proved assuming the Cassini's identity holds at rank  $n$ .

Apply the definition of  $F(n + 2)$ :

$$F(n).F(n+2) = F(n)(F(n+1)+F(n)) = F(n)^2 + F(n).F(n+1)$$

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$$F(n).F(n+2) = F(n)(F(n+1)+F(n)) = F(n)^2 + F(n).F(n+1)$$

Replace the last term using the recurrence hypothesis:

$$\begin{aligned} F(n)^2 &= F(n-1).F(n+1) - (-1)^{n+1} \\ &= F(n-1).F(n+1) + (-1)^{n+2} \end{aligned}$$

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Thus,

$$\begin{aligned} F(n).F(n+2) &= F(n).F(n+1) + F(n-1).F(n+1) + (-1)^{n+2} \\ &= F(n+1)(F(n) + F(n-1)) + (-1)^{n+2} \end{aligned}$$



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Thus,

$$\begin{aligned} F(n).F(n+2) &= F(n).F(n+1) + F(n-1).F(n+1) + (-1)^{n+2} \\ &= F(n+1)(F(n) + F(n-1)) + (-1)^{n+2} \end{aligned}$$

Apply again the definition of Fibonacci sequence

$$F(n) + F(n-1) = F(n+1), \text{ we obtain:}$$

$$F(n).F(n+2) = F(n+1)^2 + (-1)^{n+2}$$

## Computing $F(n)$ fast

$F(n)$  can be computed in  $\log_2(n)$  steps.

**Proposition.**

For all integers  $n$ :

**(a)**  $F(2n) = (F(n))^2 + (F(n-1))^2$ ;

**(b)**  $F(2n+1) = F(n) \times (2F(n-1) + F(n))$ .

## Details (a) – Proof by induction

The base case  $n = 1$  is true because

$$F(2) = (F(1))^2 + (F(0))^2 = 2$$

$$F(3) = F(1) \times (2F(0) + F(1)) = 3$$

Assume that the property holds for  $n$ , for both  $F(2n)$  and  $F(2n + 1)$ .

$$\begin{aligned} F(2(n+1)) &= F(2n+1) + F(2n) \\ &= (F(n))^2 + (F(n-1))^2 + F(n) \times (2F(n-1) + F(n)) \\ &= (F(n))^2 + (F(n-1))^2 + 2(F(n) \times F(n-1)) + (F(n))^2 \\ &= (F(n) + F(n-1))^2 + (F(n))^2 \\ &= (F(n+1))^2 + (F(n))^2 \end{aligned}$$

## Details (b)

We again start by applying the defining recurrence of the Fibonacci numbers on  $F(2(n+1) + 1)$

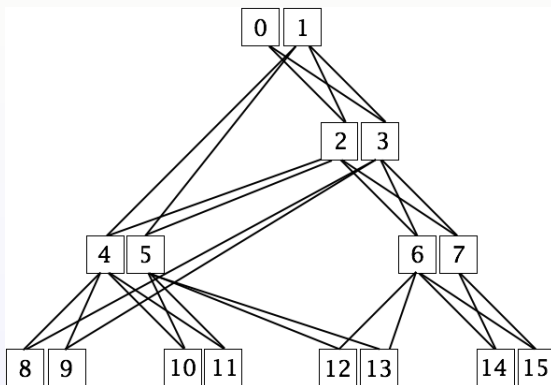
$$= F(2(n+1)) + F(2n+1)$$

$$= (F(n+1))^2 + F(n)^2 + F(n) \times (2F(n-1) + F(n))$$

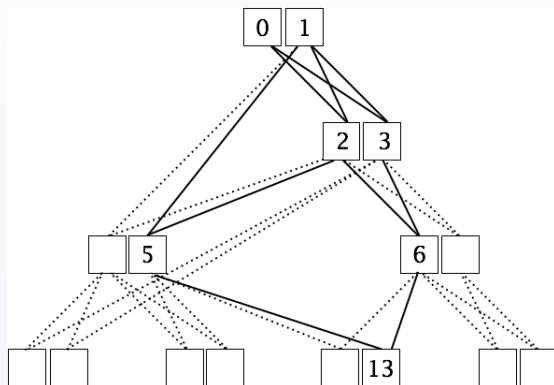
$$= (F(n+1))^2 + 2(F(n-1) + F(n)) \times F(n)$$

$$= (F(n+1))^2 + 2F(n+1) \times F(n)$$

# Pictorially



## Pictorially (from one node)



## Definition of Lucas' numbers

A natural question is:

what happens if we change the first ranks of the sequence keeping the same recurrence pattern?

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what happens if we change the first ranks of the sequence keeping the same recurrence pattern?

It has been studied by the french mathematician Edouard Lucas, starting at 2 and 1 .

For some reasons that will be clarified later, the sequence is shifted backwards (we take the convention  $L(-1) = 2$ ).



## Definition of Lucas' numbers

### Definition:

Given the two numbers  $L(0) = 1$  and  $L(1) = 3$

all the other Lucas' numbers are obtained by the same progression as Fibonacci:

- $L(n + 1) = L(n) + L(n - 1)$

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n	-1	0	1	2	3	4	5	6	7	8	9	10
F(n)		1	1	2	3	5	8	13	21	34	55	...
L(n)	2	1	3	4	7	11	18	29	47	76	123	...

- There are<sup>2</sup> strong links with Fibonacci numbers.

In particular, we established before that:

$$F(n+2) = 1 + \sum_{k=0}^n F(k).$$

We have similarly:

$$L(n+2) = 1 + \sum_{k=-1}^n L(k)$$

since the basic step of the induction is still valid<sup>3</sup>.

$$L(2) = L(-1) + L(0) + 1 = 2 + 1 + 1 = 4.$$

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<sup>2</sup>of course

<sup>3</sup>It will be true for all the progressions where  $u_1 = 1$

## A first Property

We can also easily show that the Lucas number of order  $n$  is the sum of two Fibonacci numbers:

**Proposition.**

$$L(n) = F(n - 1) + F(n + 1) \text{ for } n \geq 1$$

Let *check* this property on the first ranks:

$$n = 2, L(2) = F(1) + F(3) = 1 + 3 = 4$$

$$n = 3, L(3) = F(2) + F(4) = 2 + 5 = 7$$

$$n = 4, L(4) = F(3) + F(5) = 3 + 8 = 11$$

$$n = 5, L(5) = F(4) + F(6) = 5 + 13 = 18$$

...

## Proof by induction

- The **basis case** (for  $n = 1$ ) is true since  
 $L(1) = 3 = F(2) + F(0) = 2 + 1$ .
- **Induction step:** Let assume the property holds at all ranks  $k \leq n$  and compute  $L(n + 1)$ :  
Apply the definition of Lucas' numbers:  
 $L(n + 1) = L(n) + L(n - 1)$

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Apply the definition of Lucas' numbers:
$$L(n + 1) = L(n) + L(n - 1)$$
Apply the induction hypothesis on both terms:
$$L(n + 1) = F(n + 1) + F(n - 1) + F(n) + F(n - 2)$$

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Apply the induction hypothesis on both terms:
$$L(n + 1) = F(n + 1) + F(n - 1) + F(n) + F(n - 2)$$
Apply now the definition of Fibonacci numbers for
$$F(n + 1) + F(n) = F(n + 2) \text{ and } F(n - 1) + F(n - 2) = F(n)$$

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 Apply the induction hypothesis on both terms:  

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 Apply now the definition of Fibonacci numbers for  
 $F(n + 1) + F(n) = F(n + 2)$  and  $F(n - 1) + F(n - 2) = F(n)$   
 replace them in the previous expression:  

$$L(n + 1) = F(n + 2) + F(n)$$

which concludes the proof.



## Extension 1

Notice that using a similar approach, we obtain

$$L(n) = F(n + 2) - F(n - 2).$$

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What happens if we generalize?

**Proposition.**

$$2.L(n) = F(n+3) + F(n-3)$$

**Proof.**

We start from  $L(n) = F(n+2) - F(n-2)$

$$F(n+2) = F(n+3) - F(n+1) \text{ and}$$

$$F(n-2) = F(n-1) - F(n-3)$$

$$L(n) = F(n+3) - (F(n+1) + F(n-1)) + F(n-3)$$

$$2.L(n) = F(n+3) + F(n-3)$$

## Extension 2

Go to the next step using the same technique:

$$\begin{aligned} 2. L(n) &= F(n+3) + F(n-3) \\ &= F(n+4) - F(n+2) + F(n-2) - F(n-4) \end{aligned}$$

$$3. L(n) = F(n+4) - F(n-4)$$

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One more step:  $5. L(n) = F(n+5) + F(n-5)$

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One more step:  $5.L(n) = F(n+5) + F(n-5)$

Thus, we guess the following expression.

**Proposition<sup>4</sup>.**

$$F(k-1).L(n) = F(n+k) + (-1)^{k-1}F(n-k) \text{ for } k \leq n$$

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<sup>4</sup>The formal proof is left to the reader

## Two other propositions

Proposition.

$$F(n+1) = \frac{1}{2}(F(1).L(n) + F(n).L(1))$$

The proof comes from direct arithmetic manipulations:

$$2.F(n+1) = F(n+1) + F(n+1)$$

$$= F(n+1) + F(n) + F(n-1)$$

$$= L(n) + F(n)$$

$$= F(1).L(n) + F(n).L(1)$$

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$$= F(n+1) + F(n) + F(n-1)$$

$$= L(n) + F(n)$$

$$= F(1).L(n) + F(n).L(1)$$

The previous property can be extended for any  $k > 1$

Let compute the expression of  $F(k).L(n) + F(n).L(k)$



## A final natural question

The golden ratio.

It is a well-known result that the ratio of two consecutive Fibonacci number tends to the Golden ratio:

$$\lim_{n \rightarrow \infty} \frac{F(n)}{F(n-1)} = \Phi$$

As this result is obtained by solving the following equation  $x^2 = x + 1$  ( $\Phi$  is the positive root) and does not depend on the first rank, this holds also for the Lucas' numbers.

## A last result: the Zeckendorf's Theorem

**Objective:** Study the Fibonacci numbers as a numbering system.

Let us first introduce a notation:  $j \gg k$  iff  $j \geq k + 2$ .

The *Zeckendorf's theorem* states that:

every positive integer  $n$  has a unique decomposition of the form:  
 $n = F_{k_1} + F_{k_2} + \dots + F_{k_r}$  where  $k_1 \gg k_2 \gg \dots \gg k_r$  and  $k_r \geq 2$

Here, we assume that the Fibonacci sequence starts at index 1 and not 0, moreover, the decompositions will never consider  $F(1)$  (since  $F(1) = F(2)$ ).

## Derangements

Derangements represent one of the simplest forms of *avoidance problems*.

- A professor views it as a win-win strategy for the students in her class to grade each others' essays on *The Essential Truth in the Universe*.  
The essays thereby get graded faster.
- Moreover, each student gets a chance to see how another student has interpreted some basic component of the human experience.
- The only complication is: How should we allocate essays among the students?

*The process must ensure that no student is assigned her own essay to critique.*

This challenge is known as a *derangement problem*.

## Derangements

- A *derangement* of a (finite) set  $A$  is a *bijection*  $f : A \leftrightarrow A$  that has no *fixed point*.

In other words, for every  $a \in A$ , we must have  $f(a) \neq a$ .

Clearly, derangements always exist (for  $n > 1$ ).

One can just label the elements of set  $A$  by the numbers  $0, 1, \dots, |A| - 1$  and specify  $f(a) = a + 1 \pmod{|A|}$ .

## Playing around with a simple example

However, derangements are not so common! In fact, the set  $A = \{0, 1, 2\}$  admits six self-bijections, but only two are derangements. Which ones?

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However, derangements are not so common! In fact, the set  $A = \{0, 1, 2\}$  admits six self-bijections, but only two are derangements. Which ones?

$f(a) = a + 1 \pmod{3}$  : which maps  $(0 \rightarrow 1), (1 \rightarrow 2), (2 \rightarrow 0)$

$g(a) = a - 1 \pmod{3}$  : which maps  $(0 \rightarrow 2), (1 \rightarrow 0), (2 \rightarrow 1)$

- How many derangements does an arbitrary  $n$ -element set  $A$  have? We denote this quantity by  $d(n)$ .

# Derangements

We compute  $d(n)$  for arbitrary integer  $n$  via the following recursion:

- For  $n = 1$ :  $d(1) = 0$ .

The unique bijection in this case consists only of a fixed point.

- For  $n = 2$ :  $d(2) = 1$ .

There are two bijections in this case

- the identity, which has two fixed points
- the swap, which is a derangement.

## The inductive expression

- For  $n > 2$ :  $d(n) = (n - 1)(d(n - 1) + d(n - 2))$ :

To see this, note first that in any derangement, the first element of  $A$ , call it  $a$ , must map to some  $b \neq a$ .

- Note next that there are  $n - 1$  ways to choose  $b$ .
- There are  $d(n - 2)$  derangements under which  $b$  maps to  $a$ .  
In those cases, we know everything about  $a$  and  $b$ , so we need worry only about the remaining elements of  $A$ .  
These  $n - 2$  elements can “derange” in all possible ways.
- There are  $d(n - 1)$  derangements under which element  $b$  does not map to  $a$ .



## An observation

The preceding reasoning verifies the following recurrence

$$d(n) = \begin{cases} 0 & \text{if } n = 1 \\ 1 & \text{if } n = 2 \\ n(d(n-1) + d(n-2)) & \text{if } n > 2 \end{cases}$$

## Solving the recurrence

There are several ways to solve this recurrence.

- We can reduce the bilinearity by a linear recurrence:

$$d(n) = \begin{cases} 0 & \text{if } n = 1 \\ n \cdot d(n-1) + (-1)^n & \text{if } n > 1 \end{cases}$$

Interestingly, as the number of objects in the set to be deranged grows without bound.

The proportion of bijections that are derangements tends to the limit  $1/e$ , where  $e$  is the base of natural logarithms.

## Concluding remarks and take home message

We are more interested here by the dynamic process of *computing* the successive terms of the progressions than in their asymptotic limits when  $n$  is large...